Synchronous Dataflow Programming CS684: Embedded Systems Topic 5

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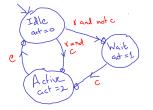
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- A node can be in exactly one mode at each clock cycle.
- Equations of the currently active mode are applied.
- Each output and internal variable has exactly one equation in each mode.
- Each mode acts as a name space and clock domain.
 All pre(x) values are stored in a mode local copy. last(x) variables are global and shared between modes.
- reset blocks can be used to reset the equations under specified conditions.

Mixed language for Multi-mode Complex Control

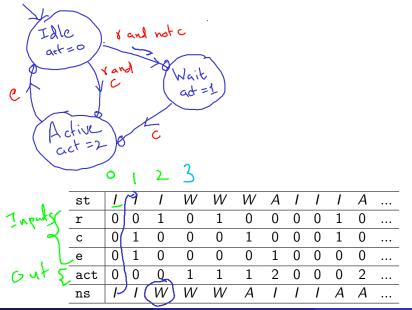
Finite State Automata with data flow equations. Hybrid Program

- States are modes.
- Each state has an associated set of equations.
- Transitions specify conditions for state (i.e. mode) change.
- Complex control is organized as automata with hierarchy, concurrency and sharing of flows.
- Dataflow and fsm control can be freely mixed and nested.



```
node myautomaton(r,c,e:bool) returns (act:int)
let
     automaton
        state Idle
                do act = 0
        until r and c then Active
            | r and not c then Wait
        state Wait
                do act = 1
        until c then Active
        state Active
                do act = 2
        until e then Idle
    end
tel
```

First Example Diagram



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Automata under Weak transitions: until

Structure of Automaton

- A set of states with transitions between them.
- One set of equations (i.e. mode block) for each state.
- A set of transitions going out of each state (keyword until or unless)
- Each transition has a guard giving condition under which it is taken and target state.

Execution of automaton: In each cycle

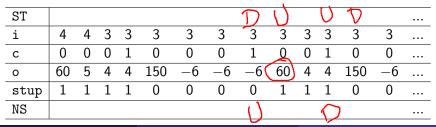
- Start State:
- With weak transitions, the start state is the Active state.
- Equations of the active state are applied.
- Guard is evaluated AFTER evaluating the active state equations.
 Guard can refer to variables defined by the equations.
 If guard true transition is taken and next state changed.
- Next State: this is the start state of the next cycle.
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State as Mode Block



Same automaton with reset transitions

```
node myautomaton(i : int; c: bool) returns (o: int; stup:bool)
let
    automaton
    state Up
    do o = 60 -> i+1; stup = true;
    until c then Down
    state Down
    do o = 150 -> -2 * i; stup = false;
    until c then Up
    end
tel
```



- A then transition resets the mode block on entry to the state,
- A continue transition enters the states mode block WITHOUT resetting.
- Each state is a mode with its own name space and clock domain.
- pre(x) in a mode refers to the previous value of x when the automaton was in this state.
- last x refers to global variable x shared between states. Its value is value of x in previour cycle (irrespective of the state).

```
node myautomaton() returns (y:int; stup:bool; v:int)
var last x:int = 2;
let
    y = x;
    automaton
        state Up
                  var w:int:
                  do x = (last x) + 1; stup = true;
                     w = 0 \rightarrow pre(w)+1; v=w;
        until x \ge 5 continue Down
        state Down
             var w:int;
             do x = (last x) - 1; stup = false;
             w = 50 \rightarrow pre(w)-2; v=w;
        until x <= 3 continue Up
    end
tel
```

```
node myautomaton() returns (y:int; stup:bool; v:int)
var last x:int = 2:
let.
    y = x;
    automaton
         state Up
                   var w:int;
                   do x = (last x) + 1; stup = true;
                       w = 0 \rightarrow pre(w)+1; v=w;
         until x \ge 5 continue Down
         state Down
             var w:int;
             do x = (last x) - 1; stup = false;
             w = 50 \implies pre(w) - 2; v = w;
         until x <= 3 continue Up
    end
tel
                                            3
                         3
                                                 4
                                                     5
                                                                3
                y
                                  5
                                                           4
                                                                     4
                                       4
                                                                          . . .
                                  1
                                            0
                                                  1
                                                      1
                                                                0
                                                                     1
                                                           0
                stup
                                                                          . . .
                                  2
                                     50
                                            48
                                                 3
                                                      4
                                                          46
                                                                44
                                                                     5
                         0
                v
                                                                          . . .
                 25
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```

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